

WHAT IS CLAIMED IS:

1                   1.       A method of authenticating a set of N information blocks, said method  
2 comprising:

3                               obtaining an initial hash value for a set of N information blocks,  
4 wherein N is an integer;

5                               altering one of said N information blocks from said set of N  
6 information blocks so as to form a revised set of N information blocks;

7                               calculating a revised hash value for said revised set of N information  
8 blocks; while

9                               calculating a check hash value for said N information blocks; then

10                              comparing said check hash value with said initial hash value;

11                              accepting said revised hash value for said revised set of N information  
12 blocks if said check hash value matches said initial hash value.

1                   2.       The method as described in claim 1 wherein said calculating said  
2 revised hash value while calculating said check hash value comprises:

3                               calculating said revised hash value in parallel with said check hash  
4 value.

1                   3.       The method as described in claim 1 wherein said calculating said  
2 revised hash value while calculating said check hash value comprises:

3                               hashing said altered block of data so as to obtain a first hashing result;

4                               storing said first hashing result in a processor; and then

5                               hashing the corresponding unaltered block of data so as to obtain a  
6 second hashing result.

1                   4.       The method as described in claim 1 wherein said calculating said  
2 revised hash value while calculating said check hash value comprises:

3 concurrently hashing said altered block of data so as to obtain a first  
4 hashing result and hashing the corresponding unaltered block of data so as to obtain a second  
5 hashing result.

1 5. The method as described in claim 1 wherein said calculating said  
2 revised hash value while calculating said check hash value comprises:

3 utilizing a single processor to calculate said revised hash value and to  
4 calculate said check hash value.

1 6. The method as described in claim 1 and further comprising:  
2 performing a linear hash of said set of data by hashing said N blocks of  
3 data in sequential order from block 1 to block N.

1 7. The method as described in claim 1 wherein said obtaining said initial  
2 hash value for said set of N information blocks comprises:

3 hashing each of said N information blocks in said set of N information  
4 blocks.

1 8. The method as described in claim 1 and further comprising:  
2 storing said initial hash value in a processor.

1 9. The method as described in claim 1 wherein said altering one of said N  
2 information blocks comprises:

3 storing a new value for at least part of one of said N information  
4 groups.

1 10. The method as described in claim 1 wherein said comparing said check  
2 hash value with said initial hash value comprises:

3 determining whether said check hash value and said initial hash value  
4 are exactly the same.

1 11. The method as described in claim 1 wherein said accepting said  
2 revised hash value comprises:

- 3 replacing said initial hash value with said revised hash value.
- 1 12. The method as described in claim 1 and further comprising:
- 2 storing the new revised hash value in the memory area previously
- 3 occupied by the initial hash value.
- 1 13. The method as described in claim 1 and further comprising:
- 2 not accepting said revised hash value as a replacement for said initial
- 3 hash value if said check hash value does not match said initial hash value.
- 1 14. The method as described in claim 13 and further comprising:
- 2 indicating a failure to authenticate.
- 1 15. The method as described in claim 1 and further comprising:
- 2 utilizing said set of data for digital rights management.
- 1 16. The method as described in claim 1 and further comprising:
- 2 replacing said initial hash value with said revised hash value.
- 1 17. The method as described in claim 1 and further comprising:
- 2 receiving as part of an initialization routine a length of a data set to be
- 3 hashed, wherein said data set is comprised of said N information groups.
- 1 18. The method as described in claim 17 and further comprising:
- 2 padding at least one of said N information groups so that each of said
- 3 N information groups is of equal length.
- 1 19. The method as described in claim 1 and further comprising:
- 2 initializing a processor so as to perform a hashing routine.
- 1 20. The method as described in claim 1 and further comprising:
- 2 initializing a hashing routine by entering the length of said set of data.

- 1                    21.     The method as described in claim 1 and further comprising:  
2                                dividing the set of data into a plurality of blocks.
- 1                    22.     The method as described in claim 1 and further comprising:  
2                                dividing the set of data into a plurality of blocks of data;  
3                                padding the last block of data so that each of said blocks of data is of  
4 equal length.
- 1                    23.     A method of authenticating a set of N information blocks, said method  
2 comprising:  
3                                obtaining an initial root key for a set of data comprised of a plurality of  
4 blocks of data, said root key operable for authenticating said set of data;  
5                                calculating hash keys for said plurality of blocks of data so that each of  
6 said hash keys corresponds to only one of said blocks of data and so that each of said blocks  
7 of data corresponds to only one of said hash keys;  
8                                storing said hash keys for said plurality of blocks of data;  
9                                altering one of said blocks of data so as to form a revised block of data;  
10                              calculating a second hash key for said revised block of data, wherein  
11 said revised block of data immediately prior to being revised corresponds to a first hash key  
12 and wherein said first hash key is one of said hash keys for said plurality of blocks of data;  
13                              utilizing said stored hash keys, including said first hash key, to  
14 calculate a check root key while utilizing said stored hash keys and said second hash key  
15 substituted in place of said first hash key to calculate a new root key;  
16                              comparing said check root key with said initial root key;  
17                              accepting said new root key if said check root key matches said initial  
18 root key.

1                   24.     The method as described in claim 23 wherein said utilizing said stored  
2 hash keys, including said first hash key, to calculate said check root key is done in parallel  
3 with said utilizing said stored hash keys and said second hash key substituted in place of said  
4 first hash key to calculate said new root key.

1                   25.     The method as described in claim 24 and further comprising:  
2                   computing a branch key;  
3                   hashing said branch key and said first hash key; and  
4                   hashing said branch key and said second hash key.

1                   26.     The method as described in claim 24 and further comprising:  
2                   computing a branch key;  
3                   hashing said branch key and said first hash key; while  
4                   hashing said branch key and said second hash key.

1                   27.     The method as described in claim 24 and further comprising:  
2                   computing a branch key; and concurrently  
3                   computing a result from said branch key and said first hash key; while  
4                   computing a result from said branch key and said second hash key.

1                   28.     The method as described in claim 24 and further comprising:  
2                   utilizing a single processor to calculate said check root key and said  
3 new root key.

1                   29.     The method as described in claim 23 and further comprising:  
2                   dividing an initial set of data into X blocks, where X is equal to 2  
3 raised to the Y power and where Y is an integer.

1                   30.     The method as described in claim 23 and further comprising:

2                               calculating intermediate branch keys by hashing previously determined  
3 branch keys; and then

4                               utilizing said intermediate branch keys to calculate said new root key.

1                   31.     The method as described in claim 23 and further comprising:  
2                               encrypting said hash keys for said plurality of blocks; and  
3                               storing said encrypted hash keys in memory outside of a processor.

1                   32.     The method as described in claim 23 and further comprising:  
2                               storing said hash keys for said plurality of blocks in a processor.

1                   33.     The method as described in claim 23 and further comprising:  
2                               storing said root key inside a processor.

1                   34.     The method as described in claim 23 wherein said altering one of said  
2 blocks of data comprises:

3                               storing a new value for at least part of one of said information groups

1                   35.     The method as described in claim 23 wherein said comparing said  
2 check root key with said initial root key comprises:

3                               determining whether said check root key and said initial root key are  
4 exactly the same.

1                   36.     The method as described in claim 23 wherein said accepting said new  
2 root key comprises replacing said initial root key with said new root key.

1                   37.     The method as described in claim 36 and further comprising:  
2                               storing said new root key in a processor in a memory area previously  
3 occupied by said initial root key.

1                   38.     The method as described in claim 23 wherein said set of N information  
2 blocks is at least partially utilized for managing digital rights.

1                    39.     The method as described in claim 23 wherein said set of N information  
2 blocks is at least partially utilized as an entitlement control message for receiving a program.

1                    40.     The method as described in claim 23 and further comprising:  
2                                initializing a hashing function by receiving the length of said N  
3 information blocks.

1                    41.     The method as described in claim 40 and further comprising:  
2                                padding the final block of the N information blocks prior to hashing  
3 the Nth block.

1                    42.     The method as described in claim 23 and further comprising:  
2                                initializing a hashing function.

1                    43.     The method as described in claim 23 and further comprising:  
2                                obtaining a set of data; and  
3                                dividing said set of data into a plurality of blocks.